Algorithms

An **algorithm** is a sequence of ordered instructions that are followed step-by-step to solve a problem. This does not need to be on a computer.

Decomposition is the breaking down of a complex problem into smaller more manageable problems that are easier to solve.

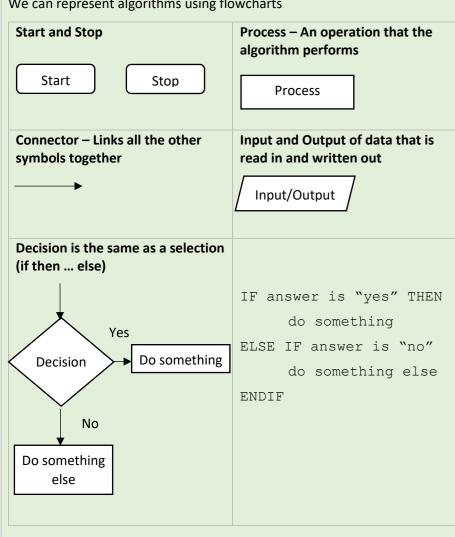
Abstraction allows us to remove unnecessary detail from a problem leaving us with only the relevant parts of a problem thereby making it easier to solve.

Algorithm Efficiency More than one algorithm can be used to solve the same problem. Normally we use the algorithm that solves the problem in the quickest time with the fewest operations or makes use of the least amount of memory.

Dry run testing is carried out using **trace tables**. The purpose of the trace tables is for the programmer to track the value of the variables and outputs at each step of the program and to track how they change throughout the running of the program.

Flowchart Symbols

We can represent algorithms using flowcharts



Pseudocode

We can represent algorithms using pseudocode

	Example	Python equivalent		
/ariable assignment	a ← 10	a = 10		
Constant assignment	constant PI ← 3.142	PI = 3.142		
nput	a ← USERINPUT	a = input()		
Output	OUTPUT "Bye"	print("Bye")		
Arithmetic Operators				
Add Multiply Divide Gubtract nteger division	+ * / - a ← 7 DIV 2	+ * / - a= 7 // 2		
Modulus (remainder)	a ← 7 MOD 2	a = 7 % 2		
Relational Operators				
Less than Greater than Equal to Not equal to Less than or equal to Greater than or equal	<pre>< > > = # or <> ≤ ></pre>	< > > == != <= >=		
Boolean Operators				
AND OR NOT	AND OR NOT	AND OR NOT		
Selection				
if	IF i > 2 THEN j ← 10 ENDIF	if i > 2: j=10		
if else	IF i > 2 THEN j ← 10 ELSE j ← 3 ENDIF	<pre>if i > 2: j=10 else: j=3</pre>		
if else if else	IF i ==2 THEN j ← 10 ELSE IF i==3 THEN	<pre>if i ==2: j=10 elif i==3: j=3</pre>		

	j ← 3	else:
	ELSE	j=1
	j ← 1	
	ENDIF	
	FNDIE	
Iteration		
While loops	a ← 1	while a<4:
	WHILE a < 4	print(a)
	OUTPUT a	a=a+1
	a ← a + 1	
	ENDWHILE	
		for a in
For loops	FOR a ← 0 TO 3	range(3):
	OUTPUT a	print(a)
	ENDFOR	1 (37)
	a ← 1	
	~ · ·	
Repeat loops	REPEAT	
Nepeut 100p3	OUTPUT a	
	a ← a + 1	
	until a←4	
Coderations	UNIIL A-4	
Subroutines		
procedure	SUB hello()	dof bollo().
	OUTPUT "hello"	<pre>def hello(): print("hello")</pre>
	ENDSUB	princ(nerro)
Function (with	FINDSOR	
paramerters and return)	OUD - 1-1 ()	1-6 -11/>
	SUB add(n)	<pre>def add(n):</pre>
	a ← 0	a=0
	FOR a ← 0 TO n	for a in
	a ← a + n	range(n+1):
	ENDFOR	a=a+n
	RETURN a	return a
	ENDSUB	
Built-in functions		
Length of array	LEN(a)	len(a)
Random integer	RANDOM_INT(0, 9)	import random
		random.randint(0,9)

Searching Algorithms

Linear Search Algorithm

- The purpose of the linear search algorithm is to find a target item within a list.
- Compares each list item one-by-one against the target until the match has been found and returns the position of the item in the list
- If all items have been checked and the search item is not in the list then the program will run through to the end of the list and return a suitable message indicating that the item is not in the list
- The algorithm runs in linear time. If n is the length of the list, then at worst the algorithm will make n comparisons. At best it will make 1 comparison and on average it will make (n+1)/2 comparisons.
- The performance of the algorithm will be improved if the target item is near the start of the list.

Example

Find the position of letter "Z" within the following list. Assume we do not have visibility of the list

Index position	0	1	2	3	4	5	6	7
Value	٧	Α	S	Z	Χ	R	Т	G

We compare it with the value in index position 0. We find that the value is "V" so we need to move on to the next index position. At index position 1 and 2 we still have not found Z. However, we get to index position 3 and we compare the target with the value and we find that they match, so the algorithm returns the index position and stops.

Pseudocode

```
i ← 0
x ← len(listOfItems)
pos ← -1
found ← False
WHILE i < x AND NOT found
IF listOfItems[i] == itemSearch THEN
found ← True
pos ← i + 1
ENDIF
i=i+1
ENDWHILE
OUTPUT pos</pre>
```

Binary Search Algorithm

- The binary search algorithm works on a sorted list by identifying the middle value in the list and comparing it with the search item.
- If the search item is smaller the mid element becomes the new high value for the search area.
- If the search item is larger the mid element becomes the low value for the search area.
- The keeps repeating until the search item is found.
- When the search item is found the index position of the item is returned.
- At each iteration the search are halved in size consequently this is an efficient algorithm.

Example: Binary search in operation to find 81



Pseudocode

 $low \leftarrow 1$

```
high ← LENGTH(arr)
mid ← (low + high) DIV 2
WHILE val ≠ arr[mid]
IF arr[mid] < val THEN
  low ← mid
ELIF arr[mid] > val THEN
  high ← mid
ENDIF
  mid ← (low + high) DIV 2
  ENDWHILE
OUTPUT mid
```

Linear search versus binary search

	Advantages	Disadvantages
Linear Search	 Very simple algorithm and easy to implement No sorting required Good for short lists 	 slow because it searchers through the whole list very inefficient for long lists
Binary Search	 much quicker than linear search, because it halves the search zone each step 	The list need to be ordered

Sorting Algorithms

Bubble Sort

- The purpose of sorting algorithms is to order an unordered list. Item can be ordered alphabetically or by number.
- Bubble sort steps through a list and compares pairs of adjacent numbers. The numbers are swapped if they are in the wrong order. For an ascending list if the left number is bigger than the right number the items are swapped otherwise the numbers are not swapped.
- The algorithm repeatedly passes through the list until no more swaps are needed.

Example

Sort the following sequence in ascending order using bubble sort: 5,3,4,1,2.

Pass	5	3	4	1	2	
1	3	5	4	1	2	Compare 5 and 3 – swap
	3	4	5	1	2	Compare 5 and 4 – swap
	3	4	1	5	2	Compare 5 and 1 – swap
	3	4	1	2	5	Compare 5 and 2 – swap; end of pass 1
Pass	3	4	1	2	5	Compare 3 and 4 – no swap
2	3	1	4	2	5	Compare 4 and 1 – swap
	3	1	2	4	5	Compare 4 and 2 – swap
	3	1	2	4	5	Compare 4 and 5 – no swap; end of pass 2
Pass	1	3	2	4	5	Compare 3 and 1 – swap
3	1	2	3	4	5	Compare 3 and 2 – swap
	1	2	3	4	5	Compare 3 and 4 – no swap
	1	2	3	4	5	Compare 4 and 5 – no swap; end of pass 3
	1	2	3	4	5	

Bubble sort Pseudocode

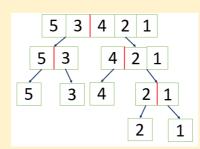
```
A=[5,3,4,1,2]
sorted ← False
WHILE not sorted
sorted ← True
FOR I TO LEN(A)-1:
IF A[i] > A[i+1]:
temp ← A[i]
A[i] ← A[i+1]
A[i+1] ← temp
sorted ← False
ENDIF
ENDFOR
ENDWHILE
```

Merge Sort

- Merge sort is a type of divide and conquer algorithm.
- There are two steps: divide and combine
- Merge sort works by dividing the unsorted list sublists. It keeps on doing this until there is 1 item in each list.
- Pairs of sublists are combined into an ordered list containing all items in the two sublists. The algorithm keeps going until there is only 1 ordered list remaining.
- Merge sort is a recursive function, that calls itself.

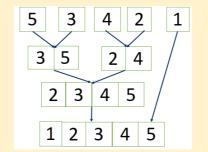
Step 1: Divide

OUTPUT A



Keep dividing until there is only 1 item in each list

Step2: Combine



- 1. The first items in the two sublists are compared, and the smallest value is copied to the parent list.
- 2. The copied item is then removed from the sublist.
- 3. When there are no items left in one of the sublists the remaining items in the other sublist are them copied in order to the parent list.

Merge sort Versus Bubble sort

	Advantages	Disadvantages
Bubble sort	Very simple and robust algorithm	Can be slow particularly for long lists. As the number of items increases the time taken for the algorithm to run increases dramatically.
Merge sort	Much faster than bubble sort especially when the number of elements is large	More complex to understand Step 1: Divide Step 2: Combine